Proposition (Prop. 3.8 [meshohi'10])
A digraph
$$\mathcal{D}$$
 on n verdiciles Containes a rooted
out-branching as a subgraph iff $rank(\mathcal{L}(\mathcal{D}))=h-1$.
In this ase, $\mathcal{N}(\mathcal{L}(\mathcal{D})) = \mathcal{A}$.

Proof: It sublices to show that "o" has algebraic multiplicity,
one iff () (intrins a routed out branching as a subgraph.
why t.
Because we know that
$$1 \in N(L(D))$$
 therefore $\operatorname{rank}(L(D)) \leq n-1$,
with equility iff o is a simple eigenvalue
Now, characteristic polynomial of $L(D)$:
 $P(N) = \operatorname{oldt}(NI - L(D))$
 $= N^n + a_{n-1} N^{n-1} + \cdots + a_1 N_1 + a_0$
where $a_{n-n} = \operatorname{sum} + \operatorname{all} \operatorname{principal} \operatorname{minors} + L(D) = \operatorname{size} k$.
 $-\operatorname{But} a_0 = \operatorname{oldt}(L(D)) = 0$.
 $-\operatorname{Thes}, \operatorname{rank} L(D) = n-1$ iff $a_1 \neq 0$.
 $-\operatorname{But}, a_1 = \sum_{v \in D} \operatorname{oldt} L_v(D)$ with removing the row
and the column corresponding to roude V.

Therefore we need to understand det
$$L_{Ve}(D)$$
.
Then [matrix-Tree theorem (undirected graph G)] =
det $L_{Ve}(G) =$ number of spanning trees in G.
Then [matrix-Tree theorem (digraph D)] =
det $L_{Ve}(D) = \sum_{T \in T_{Ve}} TT_{Ve}(e)$
 $T_{ET_{Ve}} = TT_{Ve} = Ve$

Back to the proof:
So, det
$$L_{V}(D) \neq 0$$
 if $J = V$ -rooted out branching
subgraph of D .
Thus, $a_{1} = \underset{V}{\overset{S}{=}} det L_{V}(D) \neq 0$ if $J = a$ rooted outbranching
subgraph of D .

Thus,
$$A = spm\{1\} \subseteq N(L(D))$$

Rot me in general
but it is true if D
has a no. ded out-branching
subgraph

$$\frac{Cinstant d muddin bor DAP}{V(t)^{n} = q_{1}^{T} L(D) \times (t) = 0}{\mathcal{V}(t)^{n} = q_{1}^{T} X(t)} \implies 2^{n} (t) = q_{1}^{T} \lambda(D) \times (t) = 0}{\mathcal{L} \quad b(t) \quad converges} \quad so \quad so \quad so \quad b(t) \quad when \quad closs \quad DAP \quad converges \quad to \quad de \quad average q \quad we want \quad q_{1} = 1 , \quad when \quad dees \quad des \quad de$$

Corollary: If \mathcal{D} contains a routed out-branching and is halanced, then $\mathcal{D}AP$ reaches average consensations, i.e. $\lim_{t\to\infty} \chi(t) = \frac{1^T \chi_0}{n} \cdot 1$ Proof: Recall that by hypothesis $\chi(t) \rightarrow (\mathcal{P}_1^T \chi_0) \cdot 1$ with $\mathcal{P}_1^T 1 = 1$. As \mathcal{D} is balanced, $\mathcal{P}_1 \in spon\{1\}$ $\Rightarrow \mathcal{P}_1 = \frac{1}{n} \in \mathbb{Z}$.

In fact, something more stronger is true Def: A digraph is "strongly connected" if between enery true verticies, there exists a directed path. Des: A digraph is "weakly connected" it its undirected / disoriented version is connected. Thm: The DAP on D reaches the average consonsus from every initial condition if and only if Dis weakly connected and bolanced.

Prof: = 1 if D is weakly connected and balanced
Hinework Hen it has to be strangly connected (why.t.).
However, D has a rooted out-branching subgraph.
That, because D is balanced, by the above corollary,
DAP converges to the average consumers. I.
The converges to the average consumers. I.
Conversely, suppose the convergence to average
consumers is achieved by DAP, i.e.
fin
$$X(t) = f \in UD)^{t} X(t) = \frac{1}{n} X(t)$$
. $I = \frac{1}{n} 11^{T} x_{0}$.
How every $X(t) = R^{h}$. This implies
 $\begin{bmatrix} f = e^{L(D)t} - f_{h} 12^{T} \end{bmatrix} X(t_{0}) = 0$, $\forall X(t) \in R^{h}$.
Thus, $f = e^{L(D)t} = f_{h} 11^{T}$. Now, node that
leftlight eigenvectors of $L(D)$, $e^{L(D)t}$ and $f_{h} 11^{T}$ must math,
because: $e^{L(D)t} = P e^{2(D)t} = f(L(D)t) = P J(D)P^{T}$,
and
 $f_{h} 1T = f_{h} e^{L(D)t} = P(L(D)t) = P(L(D)t) P^{T}$.
 $f_{h} to the the terms of terms of the terms of term$

Therefore, 1 has to be left mel night eigenvector of
$$L(D)$$
. By definition, $L(D) 1 = 0$. Assume $1^{T}L(D) = \alpha 1^{T}$ for some α .

But then

$$O = (L(D)I)^{T}I = I^{T}L(D)^{T}I = I^{T}(I^{T}L(D))^{T} = I^{T}(\alpha I^{T})^{T} = \alpha \cdot n$$

$$\Rightarrow \chi = O \Rightarrow I^{T}L(D) = O \Rightarrow D \text{ is balanced.}$$

Next, we have to show that D is weakly connected. Note:

$$\begin{split} \vec{e}^{\mathcal{U}(\mathcal{D})^{+}} &= \vec{P} \ \vec{e}^{\mathcal{J}(\mathcal{D})^{+}} \vec{P}^{-1} \\ &= \begin{bmatrix} 1 & n^{2} & \cdots & n^{n} \end{bmatrix} \begin{bmatrix} \vec{e}^{\mathcal{A}(\mathbf{0})^{+}} & \mathbf{0} \\ e^{-\mathcal{A}(\mathbf{0})^{+}} & \mathbf{0} \end{bmatrix} \begin{bmatrix} -\mathcal{I}_{n}^{-} - \mathbf{1}_{n}^{-} \\ -\mathcal{I}_{n}^{-} \end{bmatrix} \xrightarrow{\mathbf{1}} \mathbf{1}_{n}^{T} \\ \vec{P}_{n}^{T} &= \begin{pmatrix} -\mathcal{I}_{n}^{T} & \mathbf{1}_{n} \\ -\mathcal{I}_{n}^{T} & \mathbf{1}_{n} \end{bmatrix} \\ & \vec{P}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \\ \vec{P}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \\ & \vec{P}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T} \\ \vec{P}_{n}^{T} \mathbf{1}_{n}^{T} \mathbf{1}_{n}^{T}$$